

Comparative Analysis between Elgamal and NTRU Algorithms and their implementation of Digital Signature for Electronic Certificate

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Abstract – The emergence of electronic certificates, which are official documents in the form of digital files transmitted via the internet, facilitates the exchange of information. However, internet use has risks, such as data theft for fabricating and modifying information. This problem can be solved by applying a digital signature. The main concern in this research is how to perform a comparative analysis between asymmetric cryptographic Elgamal and NTRU (Nth-Degree Truncated Polynomial Ring) algorithms and their implementation of a digital signature as an effort to improve information security in electronic certificates. The stages of the research method are divided into the key generation process, signing, and verification. In the signing and verification process, the SHA-512 hash function is also used for hashing messages to be encrypted-decrypt and QR Code as the signature. Comparison of performance of NTRU with Elgamal algorithms required running at a pdf extension with security levels 80,128,192, 256 bits will be obtained from the templates.office.com website. The results obtained that the El Gamal algorithm is better than the NTRU algorithm, but at a higher security level, the NTRU algorithm is better than the Elgamal algorithm. In the verification experiment that has been carried out, it can be concluded that by using SHA-512 as a hash function, the N parameter used for NTRU must be greater than or equal to 512 to avoid error results from verification.

Keywords: NTRU, Elgamal, Electronic Certificate, Digital Signature, SHA-512, QR Code

1. INTRODUCTION

Electronic certificates are official documents in the form of digital files transmitted via the internet, where the internet itself is vulnerable to theft and falsification of information, such as the fabrication or modification of information [1,2,3]. To increase the security of electronic certificates. A security system in the form of digital signatures is applied to the electronic certificates. The digital signature is a cryptographic value that depends on the message's content and the message's

sender. So that different messages with the same sender will have different digital signatures [4,5,6,7].

This study aims to implement a digital signature on an electronic certificate using the NTRU and Elgamal algorithm at security levels 80,128,192, and 256 bits to see whether NTRU is better or Elgamal is better. NTRU algorithm is an asymmetric algorithm. Asymmetric algorithms have different keys during the encryption and decryption process, namely the public and private keys. The public key is the key that is published and may be known by ev-

everyone, while the private key is a key that is kept secret and may only be understood by one person [8,9]. The NTRU algorithm's security level lies in the use of polynomials during the operation process, as well as the difficulty of finding a short vector of a lattice [10,11]. The level of security of the Elgamal algorithm lies in the difficulty of calculating discrete logarithms [12]. For comparison, the Elgamal algorithm is used at security levels 80,128,192, and 256 bits to see which algorithm has better performance. Elgamal's algorithm selection is the comparison because the study [37] shows that Elgamal is a better probabilistic algorithm than RSA and has a difficulty level that lies in discrete logarithm calculations and its ability to solve fundamental distribution problems. Besides that, there are still very few studies comparing the NTRU algorithm with the algorithm Elgamal.

This research aims to implement a digital signature schemes using the NTRU and Elgamal algorithm for electronic certificates, then analyze the running time on the NTRU and Elgamal algorithms based on the process generate keys, signing, and verifying, as well as analyze the results of electronic certificate verification.

2. RELATED WORK

The NTRU algorithm is a fast and lightweight public key algorithm to provide end-to-end security that can be used to improve document security standards with better encryption and decryption than the RSA and ECC algorithms [13,14,15]. In another study, to increase document security and facilitate the validation process, the use of digital signatures using the RSA algorithm [16,17] and SHA-512 algorithms can be applied where the method used is to generate a public key and a private key with RSA. The signing process is carried out by encrypting the message digest generated from the message hashing process with SHA-512 and then verifying electronic documents by matching the results of document decryption and SHA-512 hashing of documents [18,19,20,21] in a similar study [3]. With SHA-3 hashing function and super encryption combination of RSA and AES, with QR-Code scheme to accommodate the signature code. In this implementation, the certificate will be signed with the SHA-3 hashing process sequence, encrypted with RSA, encrypted with AES, and ends by embedding the QR-Code that has been generated from the AES encryption results on the electronic certificate.

3. THEORY USED

This section discusses the theoretical background to analyze the comparison between the NTRU and Elgamal Algorithms and the implementation of digital signatures on electronic certificates.

3.1. DIGITAL SIGNATURE

Digital Signature is a means used to view authentication on digital messages, both messages transmitted through communication channels and electronic

documents; what is meant by digital signatures are signatures that have a cryptographic value that depends on the content of the message and the sender of the message, so that the message different ones with the same sender will have different digital signatures [14]. The term direct digital signature refers to a digital signature scheme that only involves the communicating party (sender, receiver). Digital signature schemes are similar to asymmetric cryptographic systems in that they involve public and private keys and run an algorithm that uses these keys to sign and verify [22,23].

3.2. NTRU (NTH-DEGREE TRUNCATED POLYNOMIAL RING UNITS) ALGORITHM

NTRU uses addition and multiplication operations in line with Ring, which is an algebraic object that has two operations, addition, and multiplication, which are related via the distributive law [22, 23], NTRU works with rings. An element will be written as a polynomial or vector according to Equation (1).

$$F = \sum_{i=0}^{N-1} F_i x^i = [F_0, F_1, \dots, F_{N-1}] \quad (1)$$

Key Generation: Choose two polynomials $f \in L_f$ $f \in$ and $g \in L_g$. Polynomial f must meet the additional requirement that it has inverse modulo q and inverse modulo p . This inverse can be expressed F_p and F_q so that the result is:

$$F_q * f \equiv 1 \pmod{q} \quad (2)$$

$$F_p * f \equiv 1 \pmod{p} \quad (3)$$

Next calculate h with Equation (4)

$$h \equiv pF_q * g \pmod{q} \quad (4)$$

Where h is a polynomial that functions as a public key and a polynomial f, fp as private keys.

NTRU encryption: Selects m as messages from a set of plaintexts L_m . Then randomly choose the polynomial L_ϕ and use the public key h to compute the encrypted message e by Equation (5)

$$e \equiv \phi * h + m \pmod{q} \quad (5)$$

Polynomial e is an encrypted message that will be sent to the recipient of the message.

NTRU decryption: In the decryption of a received e-message, the process is carried out using the private key f to calculate the value of a with Equation (6).

$$a \equiv f * e \pmod{q} \quad (6)$$

Where the coefficient a is in the interval from $q/2$ to $-q/2$. Now with a as a polynomial with integer coefficients and a private key, F_p , which can be used to recover m messages with Equation (7)

$$m \equiv F_p * a \pmod{p} \quad (7)$$

3.3. ELGAMAL ALGORITHM

Key Generation: Elgamal has a parameter of key size, which will later be used to determine positive prime numbers and integers that are primitive roots of p . To generate a public key and a private key is done by choosing a random number x , provided that than calculate the value of y with Equation (8)

$$y = g^x \pmod{p} \quad (8)$$

The result of key generation is in the form of private key x and public key y, g , and p .

Elgamal Encryption: Before performing the encryption process, first declare the message as an integer m and must lie in the range $[0, p-1]$. For large m , divide m , into smaller blocks so that each block represents a value in the range $[0, p-1]$. The encryption steps are as follows:

1. Choose a random number k , provided that $1 \leq k \leq (p - 1)$
2. Encrypt message m into value pairs (a, b) with the Equation:

$$a = g^k \pmod{p} \quad (9)$$

$$b = y^k m \pmod{p} \quad (10)$$

The pairs a and b are the ciphertext for message m . So, the ciphertext size is twice the size of the plaintext.

Elgamal's description: For decryption, the private key x is used to decrypt a and b into plaintext m with Equation [11].

$$m = b(a^x)^{-1} \pmod{p} \quad (11)$$

3.4. SHA-512 (SECURE HASHING ALGORITHM)

SHA (Secure Hashing Algorithm) is designed by the National Security Agency (NSA). SHA security is based on the fact that a birthday attack on a digest of n bits results in a collision with a work factor of about $2n/2$ [24]. SHA-512 is one of the results of the revision of the FIPS standard in 2002, which defines three new versions of SHA, with hash values of 256, 384, and 512-bits long, known as SHA-256, SHA-384, and SHA-512. Collectively, these hash algorithms are known as SHA-2. This new version has the same basic structure and uses the same types of binary logical operations and modular arithmetic as SHA-1 [12].

3.5. QR CODE

QR Code is a two-dimensional matrix symbology with a position detection pattern at three angles initially designed for very high-speed reading and Omnidirectional reading. The QR Code was developed to increase the speed of reading complex structured 2D barcodes. Other QR Code features are bulk data capacity, high data density, and selectable levels of error cor-

rection capability [25, 26,27,28]. QR codes store data using a graphical representation. The essence of this representation is based on the arrangement of several simple geometric shapes on a fixed space [29,30,31,32].

4. PROPOSED WORK

This section provided an overview of our solution comparative analysis of the Elgamal and NTRU Algorithm, and the Implementation of digital signatures on an electronic certificate.

Fig.1 shows the comparative analysis model of the digital signature in the NTRU and Elgamal algorithms in carrying out the signing and verifying process. The comparative analysis model consists of two processes: (1) the signing process using both NTRU and Elgamal Algorithms and (2) the verifying process using both NTRU and Elgamal Algorithm. The process begins with the user who selects the electronic certificate with a pdf extension file. This file can be called plaintext.

- a) The signing process using both NTRU and Elgamal Algorithms

The signing process begins with (1) calculating the hash value of the certificate file with the SHA-512 hash function, which produces a message digest (m). (2) these results are then encrypted with the Private Key and produce a message digest cipher (3) The cipher message digest is then stored in the form of a file pythons. (4) Generate a QR Code with the cipher message digest file address as the data. QR Code generation is done using a QR Code Generator. (5) Embed the QR Code as a signature into the certificate file. (6) Sending certificate file + QR Code to recipient.

- b) The In the verifying process using both NTRU and Elgamal Algorithm

The verifying process is executed with (1) Look for the signature in the form of a QR Code, which is contained in the certificate file, then separate the QR Code from the certificate. (2) Calculate the hash value of the certificate file, and generate a message digest (m'). (3) Decode the QR Code to get the data in it, which is the address of the cipher message digests file. (4) Based on the address obtained, then look for the cipher message digest file to get the cipher message digest. (5) Decrypt the cipher message digest with the Public Key, which results in a message digest (m). (6) Comparing m and m' . If $m = m'$, then it can be concluded that the certificate file is "Valid". Meanwhile, if it is not the same, it can be concluded that the certificate file is "Invalid"

Testing is done by running the system to see if the system has running according to its function or not, by doing signing and verifying experiments document. The certificate document will be subject to a signing process and generate a new certificate document complete with a digital signature in it. This document will then be subject to two treatments, (1) The certificate document does not subject to any content changes, so

when the verification process is carried out with the system, it will display a result indicating that the certificate is correct or valid. (2) The certificate document is subject to content changes so that when the verification process is carried out with the system, it will display the

following results indicating that the certificate is incorrect or invalid. After testing, system test results data in the form of running time during the key generation, sign, and verify processes. The data is then analyzed to see whether NTRU is better or ElGamal is better.

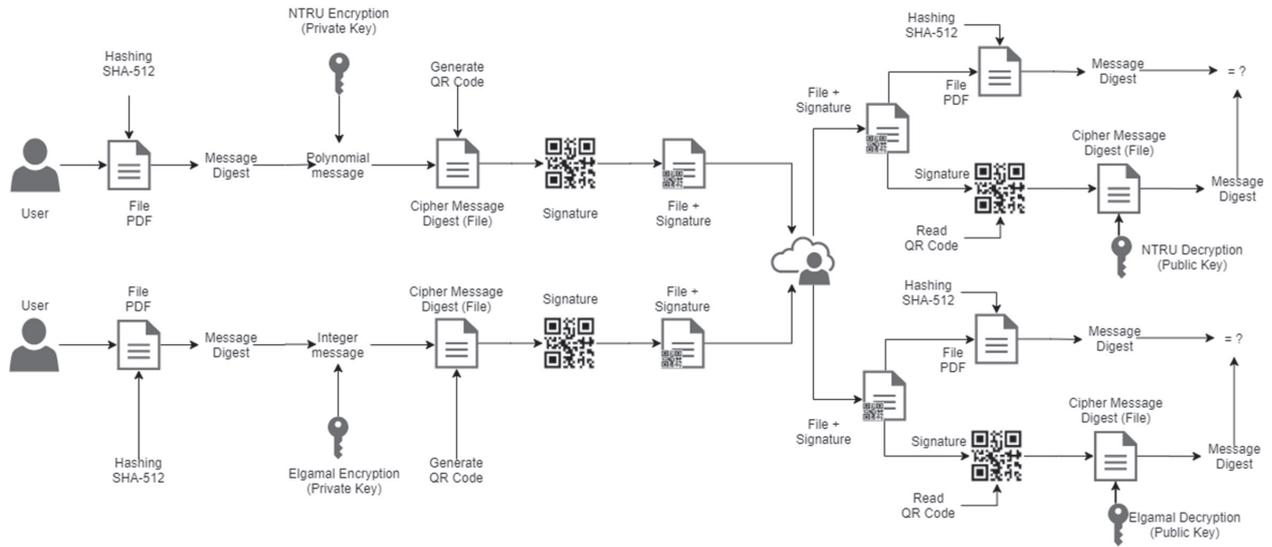


Fig. 1. Comparative Analysis of ElGamal and NTRU Algorithm and Implementation in the electronic certificate

5. RESULT AND DISCUSSION

In this section, perform a comparative analysis of the NTRU and ElGamal Algorithm and an Implementation in the electronic certificate.

5.1. DATA COLLECTION

The data for this research is a certificate with a pdf extension obtained from the templates.office.com website. The certificate can be depicted in Fig 2.

The performance of NTRU with ElGamal required data showing that the NTRU with such parameters will be comparable to ElGamal with so many bits. The data which can be used for this purpose is the security level of the NTRU algorithm with the ElGamal algorithm [33, 34, 35, 36]. The data can be seen in Table 1.

Table 1. Security Level Elgamal and NTRU

Security Level (bits)	NTRU	Elgamal (bits)
80	251	1024
128	397	3072
192	587	7680
256	787	15360

5.2. NTRU IMPLEMENTATION OF THE DIGITAL SIGNATURE SCHEME

NTRU Key Generation: The main parameters of the NTRU algorithm are integers N , p , and q , and the four sets L_f, L_g, L_m, L_ϕ polynomial of degree $N-1$ with integer coefficient. The integers p and q do not have to be prime, but provided that $\gcd(p, q) = 1$, and q will always be much greater than p . NTRU works with rings polynomial $R = \mathbb{Z}[X]/(X^N - 1)$.

For example, the low parameter is chosen to facilitate the ease of its calculation. The selected parameters are $N = 251, p = 3$, and $q = 2048$. Then randomly determine the polynomials f and g as follows:

$$g = x^{*247} + x^{*245} - x^{*244} + x^{*240} + x^{*235} + x^{*225} + x^{*221} + x^{*220} + x^{*219} - x^{*217} + \dots - x^{*36} - x^{*34} + x^{*31} - x^{*30} + x^{*29} + x^{*27} + x^{*21} + x^{*19} - x^{*16} + x^{*15} - x^{*13} - x^{*7} - x^{*2} - x - 1$$

$$f = x^{*250} + x^{*249} - x^{*248} - x^{*246} + x^{*244} - x^{*243} - x^{*242} + x^{*241} - x^{*239} + x^{*238} - x^{*237} - \dots - x^{*17} + x^{*16} - x^{*15} + x^{*14} - x^{*13} - x^{*12} - x^{*9} - x^{*8} + x^{*6} + x^{*5} + x^{*4} - x^{*3} + x^{*2} - 1$$

Then calculate the inverse of $f \text{ mod } p$ and $f \text{ mod } q$. The results obtained are

$$f_p^{-1} = x^{*250} + x^{*249} + x^{*247} + x^{*245} + x^{*244} + x^{*243} - x^{*241} - x^{*240} - x^{*239} + x^{*236} - x^{*235} + x^{*234} - \dots + x^{*22} + x^{*21} + x^{*19} - x^{*18} + x^{*17} + x^{*15} - x^{*14} + x^{*10} + x^{*8} + x^{*7} + x^{*6} - x^{*5} + x^{*4} - x^{*3} - x$$

$$f_q^{-1} = -919x^{*250} - 141x^{*249} + 376x^{*248} + 556x^{*247} + 275x^{*246} + 150x^{*245} + 201x^{*244} + \dots + 249x^{*9} - 199x^{*8} + 805x^{*7} + 384x^{*6} + 216x^{*5} + 864x^{*4} + 819x^{*3} - 696x^{*2} + 686x + 1019$$



Fig. 2. Certificate Data with a pdf extension

The last step is the public key calculation

$$h = 1013x^{250} + 684x^{249} + 737x^{248} + 355x^{247} - 113x^{246} - 991x^{245} - 652x^{244} + 473x^{243} - \dots - 571x^8 - 531x^7 - 414x^6 - 466x^5 + 847x^4 - 214x^3 - 926x^2 - 1014x + 718$$

The result of generating the NTRU key is stored in a python file, as shown in Fig. 3, and is used in the signing and verifying process as input.



Fig. 3. NTRU Key File

NTRU Signing: The signing process requires the input of a certificate file with extension pdf and a key file with extension npz. The message in the certificate file will be hashed with the SHA-512 hash function and produce a message digest. These results are then encrypted using the NTRU algorithm with completion $e \equiv r * h + m; (mod q)$ so that the results are in the form of ciphertext. The resulting ciphertext will be stored in a file and the address of the file is then used as data in QR Code generation, by utilizing the QR Code module from Python and acts as a signature.

To be clear, assume that the message is already in the form of polynomial.

$$m = x^{509} + x^{507} + x^{506} + x^{505} + x^{504} + x^{503} + x^{502} + x^{499} + x^{498} + x^{493} + x^{490} + x^{489} + \dots + x^{38} + x^{36} + x^{33} + x^{31} + x^{28} + x^{26} + x^{14} + x^{12} + x^{11} + x^{10} + x^7 + x^6 + x^5 + 2x + 2$$

Then calculate a random polynomial r of degree 251. Assume that the result is:

$$r = x^{202} + x^{186} + x^{183} + x^{180} + x^{176} + x^{169} + x^{167} + x^{163} + x^{151} + x^{146} + x^{135} + x^{127} + x^{126} + x^{123} + x^{121} + x^{114} + x^{112} + x^{109} + x^{107} + x^{100} + x^{97} + x^{94} + x^{93} + x^{92} + x^{88} + x^{84} + x^{83} + x^{82} + x^{80} + x^{75} + x^{67} + x^{66} + x^{64} + x^{60} + x^{59} + x^{58} + x^{50} + x^{46} + x^{45} + x^{38} + x^{37} + x^{36} + x^{34} + x^{33} + x^{32} + x^{31} + x^{30} + x^{29} + x^{28} + x^{27} + x^{26} + x^{25} + x^{24} + x^{23} + x^{22} + x^{21} + x^{20} + x^{19} + x^{18} + x^{17} + x^{16} + x^{15} + x^{14} + x^{13} + x^{12} + x^{11} + x^{10} + x^9 + x^8 + x^7 + x^6 + x^5 + x^4 + x^3 + x^2 + x + 1$$

Based on the encryption formula, the encrypted message value is obtained with the public key in the previous example.

$$e = -39x^{250} + 981x^{249} + 124x^{248} - 90x^{247} + 238x^{246} - 129x^{245} - 147x^{244} - 217x^{243} - \dots - 173x^{99} + 808x^{98} - 981x^{97} - 4x^{96} + 952x^{95} + 962x^{94} + 472x^{93} + 935x^{92} + 185x^9 - 946$$

The final result of the signature process (with parameters $N = 587, p = 3,$ and $q = 2048$) is a pdf certificate file, as shown in Fig 4.



Fig. 4. Signed Certificate File

During the signing process, the data, which was originally a plaintext message, was processed to become a signature. An example of the process of changing the data can be seen in Table 2.

Table 2. Data Changes During NTRU Signing

Process	Results
Read the contents of the certificate file	CERTIFICATE OF TRAINING This certifies that Tengiz Kharatishvili has successfully completed training in video publishing Rowan Murphy, Sr. Videographer June 04, 20XX Date
Hashing messages with SHA – 512	b0f741a08c914065d146114e6f946b 50449ca05a87753cf014cb572eec68 dbb5172457fa9e49a9f188a1a377ed 9047eb1be64d1d61c27f4327fc56f8 9d85fba4
Converting message to polynomial form	Poly($x^{511} + x^{510} + x^{509} + x^{508} + \dots + x^{22} + 2^*x + 2, x, \text{domain}='ZZ'$)
Encryption result	Poly($38^*x^{586} - 28^*x^{585} - \dots - 26^*x^{33} - 38^*x^{32} - 29^*x + 9, x, \text{domain}='ZZ'$)
Save the encryption result into file	encrypted_220512_191310.npz

Generate QR-Code



Verifying NTRU: The verification process is initiated by inputting the signed certificate file and key file. The certificate file will then be carried out by two different processes, namely (a) the process of hashing the contents of the certificate file, which produces a message digest, and (b) the process of decrypting the ciphertext using the completion of the NTRU algorithm $d \equiv f_p^{-1} * [f * e]_q (mod p)$. The decryption process can be done by first decoding the QR Code to obtain the ciphertext file address.

More specifically, take the values of $e, f,$ and $fp-1$ using the previous encryption calculations. Then use the private key f to calculate the value of d with the formula until you get the result:

$$d = x^{250} + x^{248} + x^{247} + x^{244} + x^{242} + x^{241} + x^{240} + x^{239} + x^{238} + x^{236} + x^{234} + x^{233} + \dots + x^{22} + x^{21} + x^{17} + x^{14} + x^{13} + x^{12} + x^{11} + x^{99} + x^{98} + x^{96} + x^{95} + x^{94} + x^{93} + x^{92}$$

An example of the process of changing the data can be seen in Table 3.

Table 3. Data Changes During NTRU Verifying

Process	File	Signature
Processed data	CERTIFICATE OF TRAINING This certifies that Tengiz Kharatishvili has successfully completed training in video publishing Rowan Murphy, Sr. Videographer June 04, 20XX Date	 QR Code Decode Result: encrypted_220512_191310.npz

	Message Digest:	Decryption Result:
	b0f741a08c914065d	b0f741a08c914065
	146114e6f946b5044	d146114e6f946b50
	9ca05a87753cf014cb	449ca05a87753cf01
Results	572eec68dbb517245	4cb572eec68dbb51
	7fa9e49a9f188a1a37	72457fa9e49a9f188
	7ed9047eb1be64d1	a1a377ed9047eb1b
	d61c27f4327cf56f89	e64d1d61c27f4327f
	d85fba4	c56f89d85fba4

Verification result	Valid Certificate
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5.3. ELGAMAL IMPLEMENTATION OF THE DIGITAL SIGNATURE SCHEME

Elgamal Key Generation: Elgamal's algorithm uses a parameter in the form of a key size, which will later be used to determine the positive prime number p and the integer q , which is the primitive root of p . More specifically, for example, a low parameter will be chosen so that calculations can be carried out easily. The selected parameter is a key length of 1024 bits.

Determine the positive prime number p and the primitive root integer (of p) g , as follows:

p :

```
12489857221665811115722087408354707200904245054949915548595594703477566992936482827403733
844778047412215234604025630607003873323702143405841378919941167032148636454799401302956338
429661580291454927374628123190632069773463945113846463535827792773210978304435023477920819
1301968003951971300258281822235019146787
```

g :

```
634187109953871722414031920698055381971645540161670912639191994243896878671751031164050423
516381048310300668121586108891709476922650706546659951979204975000220138786820104549193674
989313904985155285516103341031838636861024714927332979007191006009952392275449498672360104
06029917652700377892076852541591702248
```

After that, a random number x that meets the conditions can be determined:

x :

```
91491006234848614245179977624682077925335183449227099755560879669520524960131412926704708
34735786780210541158300926247902230874565092117625576634432767085655123007686017958418385
032352969005034382561312282392552388261175928361869315037599080053753237342233076253295610
3199301369876073313607616134516448307
```

Then calculate the value of the public key y using the formula until it gets the result:

y :

```
6513528886376233423156589926619645695596911824928693298285511568454874973191493243304338
387323226204977601743100432410851026706922059501650769220930864413183397072789375870617976
6387255086140429860544959324431905194105832001038932251989629771258967253252343837800990
05128460592394696256845127484968882332
```

The result of the Elgamal key generation is stored in a python file, as shown in Fig. 5, and is used in the signing and verifying process as input.



Fig. 5. Key File

Elgamal Signing: The signing process requires the input of a certificate file with a pdf extension and a key file with npz extension. Messages in the certificate file are hashed with SHA-512 and generate a message digest. These results are then encrypted using the Elgamal algorithm with completion $a = g^k \pmod{p}$ and $b = y^k m \pmod{p}$.

So that the results are in the form of ciphertext. The resulting ciphertext will be stored in a file and the address of the file is then used as data in QR Code generation by utilizing the QRcode module from Python and acts as a signature.

More specifically, for example, assume that the message you want to send is already in the form of an integer:

m :

```
[14538442250038144105231320892152646169614843557636457374901122892002070321499286938540856
146613390051004182394475048657577252138832684349648204124476643815169212696928541858932923
96867705009816910287787229500894085499216436358018904406817756050209176272038926760434163
5003563025971688376388564955828518655,
277051609433009417896108549721465843810604790226550854273301672281256388876166909189707041
387833196283598744810191333262733937741703000683156540838356623441981685851286315026748872
7224262036602839091795881935679447740902122928454378590798578870757748432673318501925868
63442741449112785962979225229952101, 51]
```

Then determine the random number k :

k :

```
413595623204417673776342832888503224125472768031594611495484770241284420055844750071727967
000160416200371528602432811785808733490802274781083247710926862789684457461269585324098123
751015437125155307002915256627810925348141437771779326514895333116231344373179558851453931
72700788991929703414529595208367799321
```

By using the values of g , p , and y obtained in the previous generate key, then calculating the values of a and b with the formula until the following results are obtained.

a :

```
704798831812903730767592840850362047259304136001732045823020489237667079342633510409537862
74254348456281394440796452980960199763720656329447552526439670077386342709719064201262707
090770638523139374257056105282585107662663863042102239634363113079173711537825725359070539
66206171964146582994168138422591559688
```

b :

```
381469081370942410420486521225778003348021939250172661701837602867678272021630372568428
102398400480159146431599276831352445023784702785633990400058547222602018490187243485771195
827375432340619371181532871649931584390177232329400813336145567897002704448237950502335543
92401223836825760672707159682353497160
```

The final result of the signature process is a pdf certificate file, as shown in Fig. 6.



Fig. 6. Signed (Elgamal) Certificate File

During the signing process, the data, which was initially a plaintext message, was processed to become a signature. An example of the process of changing the data can be seen in Table 4.

Table 4. Data Changes During Elgamal Signing

Process	Results
Read the contents of the certificate file	CERTIFICATE OF TRAINING This certifies that Tengiz Kharatishvili has successfully completed training in video publishing Rowan Murphy, Sr. Videographer June 04, 20XX Date
Hashing messages with SHA – 512	b0f741a08c914065d146114e6f946b50449ca05a87753cf014cb572ecc68dbb5172457fa9e49a9f188a1a377ed9047eb1be64d1d61c27f4327fc56f89d85fba4
Convert message digest to integer	2688241...950207, 2660811...905845, 52
Determine the random number k	6791467875486908727...72474959180919163072209829196994526
Encryption result	82875883331547...83931630431734, 56002328522083...40235165644727, 79852418918538...58497092617061, 31666415267698...84650783518716, 11618727619395...39185616194775, 37640656195438...72911613317981
Save the encryption result into file	keyElgamal_220512_203003.npz
Generate QR-Code	

Verifying Elgamal: The verification process begins with inputting a signed certificate file and a key file. The certificate file will then be carried out by two different processes, namely (a) the process of hashing the contents of the certificate file, which produces a message digest, and (b) the process of decrypting the ciphertext using the completion of the NTRU algorithm $m = b(a^x)^{-1} \pmod p$. The decryption process can be done by first decoding the QR Code to obtain the ciphertext file address.

More specifically, use the private keys x , ciphertext $[a, b]$, and p from the previous calculation, then calculate the value of m with the formula until the results are obtained

m:
145384422500381441052313208921526461696149435576364573749011228920020703214992869385408561
466133900510041823944750486575772521388326843496482041244766438151692126969285418589329239
68677050098169102877872295008940854992164363580189044068177560502091762720389267604341635
803563025971688376388564955828518655,
277051609433009417896108549721465843810604790226550854273301672281256388876166909189707041
387833196283598744810191333262733937741703000683156540838356623441981685851286315026748872
72242620366028390917958819356794477409021229284543785907985788077577484326733185079125868
634427414491127859629792255229952101

An example of the process of changing the data can be seen in Table 5.

Table 5. Data Changes During Elgamal Verifying

Process	File	Signature
Processed data	CERTIFICATE OF TRAINING This certifies that Tengiz Kharatishvili has successfully completed training in video publishing	

Processed data	Rowan Murphy, Sr. Videographer June 04, 20XX Date	QR Code Decode Result: keyElgamal_220512_203003.npz
Results	Message Digest: b0f741a08c914065d146114e6f946b50449ca05a87753cf014cb572ecc68dbb5172457fa9e49a9f188a1a377ed9047eb1be64d1d61c27f4327fc56f89d85fba4	Decryption Result: b0f741a08c914065d146114e6f946b50449ca05a87753cf014cb572ecc68dbb5172457fa9e49a9f188a1a377ed9047eb1be64d1d61c27f4327fc56f89d85fba4
Verification result	Valid Certificate	

5.4. COMPARISON OF NTRU AND ELGAMAL ALGORITHM

Key Generation: Longer keys will provide higher security but consume more computational time, so the value of safety and speed will be inversely related. Table 6 shows the running time results for key generation on the Elgamal and NTRU algorithms.

Table 6. Running Time of NTRU and Elgamal Key Generation

Security Level	Algorithm	Trial (seconds)					Average
		1	2	3	4	5	
Low	NTRU	97.977	95.85	95.599	95.867	114.09	99.8766
	Elgamal	5.913	1.324	2.698	25.212	119.116	30.8526
Standard	NTRU	274.4	282.074	275.681	250.869	252.186	267.042
	Elgamal	133.610	4143.273	384.657	1712.502	282.778	1331.364
High	NTRU	532.952	527.532	539.384	684.279	532.874	563.4042
	Elgamal	3875.925	6818.925	7651.898	2647.856	2641.153	4727.151
Highest	NTRU	1255.205	1154.33	1106.095	1024.373	1026.911	1113.3828
	Elgamal	>9 Hours (32400)	> 9 Hours (32400)				

From the Table 6, it can be described in graphical form, as shown in Fig. 7.

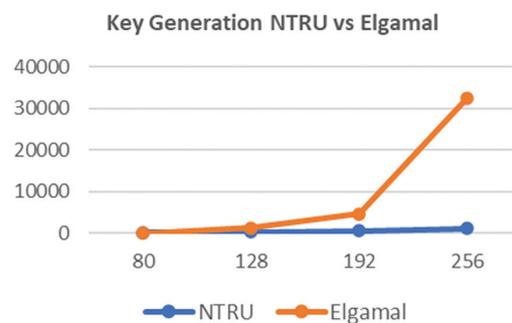


Fig. 7. Running Time of NTRU and Elgamal Key Generation

From Fig. 7, the computation time required to generate an NTRU key at a low-security level is 3x slower than Elgamal, but at a standard and high NTRU security level it is almost 5x and 8x faster than Elgamal

Signs: The time required to sign the file using the two algorithms is compared to evaluate the performance of the proposed system. The running time of the signing process with the NTRU and Elgamal Algorithms in five trials, results are shown in Table 7.

Table 7. Running Time Signing NTRU and Elgamal

Security Level	Algoritma	Trial (seconds)					Average
		1	2	3	4	5	
Low	NTRU	6.969	6.242	6.138	5.967	6.242	6.3116
	Elgamal	0.62	0.606	0.649	0.565	0.527	0.5934
Standard	NTRU	11.975	12.004	11.829	10.601	10.353	11.3524
	Elgamal	2.505	2.281	2.226	2.218	2.123	2.2706
High	NTRU	20.529	20.401	20.5	24.45	20.725	21.321
	Elgamal	28.23	27.072	31.958	34.558	31.604	30.6844
Highest	NTRU	50.471	45.614	46.38	40.799	40.158	44.6844
	Elgamal	-	-	-	-	-	-

Table 7 shows graphs of running time for the signing process using the Elgamal and NTRU algorithms, as shown in Fig. 8. Experimental running time on Elgamal with the highest security cannot be carried out due to a failure in the key generation process which cannot generate a public key and a private key, so the process cannot be continued.

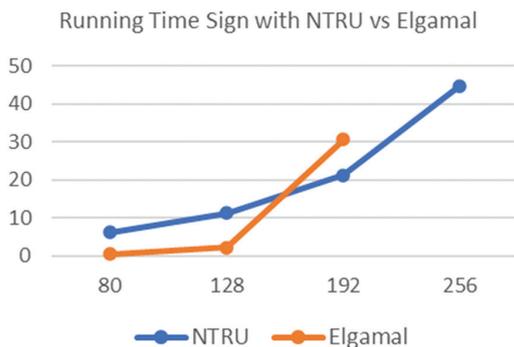


Fig. 8. Running Time Signing NTRU and Elgamal

From Fig. 8, the time for signing using Elgamal is faster than NTRU when the experiment is carried out at a low-security level and standard security. At the same time, at a higher security level, namely high safety, and highest security, the signing process using the NTRU algorithm requires faster time if compared to Elgamal's algorithm. So, at a higher level of protection, it can be said that the signing process using the NTRU algorithm is faster and safer than using the Elgamal algorithm. From Fig 8, the computation time required to sign NTRU at a high-security level is almost 1.5x faster than Elgamal.

Verify: The NTRU cryptosystem significantly produces faster average speeds than Elgamal when the key size is increased. The time required to sign the file using the two algorithms is compared to evaluate the performance of the proposed system. Table 8 shows the results of running time verifying NTRU and Elgamal

Table 8. Running Time Verifying NTRU and Elgamal

Security Level	Algoritma	Trial (seconds)					Average
		1	2	3	4	5	
Low	NTRU	8.578	7.815	7.538	7.619	8.121	7.9342
	Elgamal	0.527	0.648	0.555	0.54	0.531	0.5602
Standard	NTRU	20.843	20.985	22.962	19.06	18.214	20.4128
	Elgamal	3.614	3.610	3.458	3.443	3.509	3.5268
High	NTRU	39.782	39.462	40.5	50.229	40.109	42.0164
	Elgamal	64.338	70.588	59.873	61.728	61.132	63.5318
Highest	NTRU	101.122	83.753	84.22	77.998	79.882	85.395
	Elgamal	-	-	-	-	-	-

Table 8 shows graphs of running time Verifying NTRU and Elgamal Cryptosystem, as shown in Fig. 9.

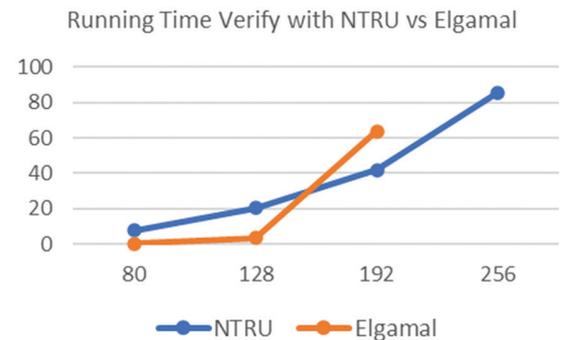


Fig. 9. Running Time Verifying NTRU and Elgamal

From Fig. 9, the verification process with the Elgamal algorithm for low-security levels, Elgamal is superior to NTRU. However, the security level is at a higher security level the NTRU is found to be faster than the Elgamal algorithm.

5.5. DIGITAL SCHEMATIC TESTING

NTRU: The test is carried out by running a digital signature scheme using the NTRU algorithm, against the same certificate file, with the final result as an "Invalid" or "Valid" certificate statement. It can be shown in Table 9.

Elgamal: The test is carried out by running a digital signature scheme using the Elgamal algorithm, against the same certificate file, with the final result in the form of an "Invalid" or "Valid" certificate statement. . It can be shown in Table 9.

Table 9 shows that in the Elgamal scheme, the verification results on certificates without changes at all security levels have "Valid" verification results, and certificates with changes have "Invalid" results. In the NTRU scheme, the verification results on the certificate with changes in all N-parameter tests result "Invalid", this happens because a difference in the contents of the certificate will result from a new message digest m (from the hashing process) which during the comparison process the value of m and m' (message digest from the decryption process) will be different and cause the verification results to be invalid.

Table 9. Testing the Digital Signature Scheme with NTRU

Security level (bits)	NTRU			Elgamal		
	N	Certificate Verification Results (Without Changes)	Certificate Verification Results (With Changes)	bits	Certificate Verification Results (Without Changes)	Certificate Verification Results (With Changes)
80	251	Invalid	Invalid	1024	Valid	Invalid
128	397	Invalid	Invalid	3071	Valid	Invalid
192	587	Valid	Invalid	7680	Valid	Invalid
256	787	Valid	Invalid	15360	-	-

However, Table 9 shows that the results of certificate verification without changes with parameter values NTRU N – 587 and N-787 show “Valid” results. The verification results do not change because the length of the polynomial ring R can include the size of the original message polynomial with the highest degree of 511. When decrypted, the decrypted polynomial will not be truncated because of the polynomial ring rule R . So that from the beginning to the end of the decryption process, the polynomial length of the processed message will not be truncated and intact. Meanwhile, the N – 251 and N -397 tests show “Invalid” results, which can occur due to the use of the SHA-512 hash function and the N parameter value that affects the length of the polynomial ring $R = \mathbb{Z}[X]/(X^N - 1)$. When the original message is hashed with SHA-512, which is then converted into polynomial form, the resulting polynomial will have a maximum length of 511 degrees. Meanwhile, during the encryption and decryption process, the size of the polynomial will follow the rules $R = \mathbb{Z}[X]/(X^N - 1)$. If the N parameter value used is 251, then the highest degree of the applicable polynomial is N-1 or 250. When the original message polynomial has the highest degree of 511 while the message decryption polynomial only has the highest degree of 250 because it follows the ring polynomial R . Then the message digest generated from the decryption process is different from the message digest from hashing the original message. So that the verification results show the same "Invalid" results, namely "Invalid" for the same reason.

6. CONCLUSION

From the problems encountered, the proposed problem-solving solutions, as well as the experiments carried out. It can be concluded that with the application of the NTRU and Elgamal Algorithms in the digital signature scheme, based on the comparison of running time in the key generation, sign, and verify processes, it shows that when the security level is low, NTRU is slower than Elgamal, but at the high-security level, NTRU is faster than Elgamal, which is 1.4x faster in the signing process and 1.5x faster in the verification process. So it can be said that NTRU, at a higher level of security, has faster when compared to Elgamal. In addition, the results of the certificate verification test with NTRU and Elgamal on the digital signature have been tested to be safe. This is proven by the testing process where data

modification is carried out in the certificate document, and the program manages to find out and shows the results "Invalid ". In the certificate document without modification, the program shows the result "Valid", but this result does not apply to NTRU N-251 and NTRU N-397 and shows the result "Invalid," which should be "Valid" this can happen because of the role of bit length The SHA used is SHA-512.

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